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Performance Analysis of Cellular Radio Access Networks Relying on Control- and User-Plane Separation

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Abstract—In a cellular radio access network (RAN) relying on control- and user-plane separation, control base stations (CBSs) form the control plane (CP) and take care of control coverage as well as low speed data delivery. In contrast, traffic base stations (TBSs) in the user plane (UP) are used for transmitting high speed data. Based on queuing theory and stochastic geometry, we derive the coverage probability of the CP under non-line-of-sight (NLoS) transmissions, and derive the energy efficiency of the UP under coexisting line-of-sight (LoS) and NLoS propagation modes. Numerical results show that when the traffic load is high, the coverage probability of the CP increases with the increasing CBSs' data rate or the density of CBSs. We also found that the maximum energy efficiency of the UP can be achieved by optimizing the TBS density, while achieving a high area spectral efficiency.

Index Terms—Cellular RANs, control- and user-plane separation, coverage probability, energy efficiency, LoS, NLoS, queuing theory

I. INTRODUCTION

In response to the enormous increase in mobile data traffic, cellular radio access networks (RANs) relying on control- and user-plane separation (CUPS) have been introduced in [1]. CUPS decouples the control coverage from traffic coverage, by which the traffic coverage may be adapted in response to both the dynamics of mobile traffic characteristics and to the prevalent quality of service (QoS) requirements. A cellular RAN relying on CUPS consists of two kinds of base stations (BSs). Specifically, the control base stations (CBSs) of the control plane (CP) deliver control signals (such as cell-specific reference signals and radio resource control, see [2] for details) and low rate signals at a low carrier frequency, while traffic

base stations (TBSs) in the user plane (UP) provide high-rate data transmissions at a high carrier frequency.

The idea of physical separation the CP and the UP is closely in line with 5G standards in terms of Non-Standalone (NSA) (3GPP TR 38.801). In the NSA mode, the BSs are classified into two types that are in charge of control and information transmissions, respectively. CUPS achieves the advantages of reduced design complexity and expenditure of BS, simplified network management through software defined networking (SDN) [3] and reduced energy consumption at TBSs [1]. Moreover, the idea of CUPS will be used in 5G/6G as the key enabling technology of open network architecture [4], which facilitates the network with the characteristics of fast-development, flexibility, re-configurability and economy.

The system model of cellular RANs relying on CUPS studied in [5]–[7] assumes that the fading channel models of both planes are the same and that only non-line-of-sight (NLoS) transmissions exist. In [5] and [6], the authors investigated caching strategies of cache-enabled hyper-cellular networks (HCNs) based on the CUPS. Zhang *et al.* [7] studied the benefits of the sleep scheme of small cells in HCNs, where the densities of macro BS users and small cell users are proportional to the traffic load in each layer. The channel models used in aforementioned literatures only considered the NLoS path loss mode. However, the shorter transmission distance in the UP results in predominantly line-of-sight (LoS) transmissions [8]. In [9], the authors studied CUPS based RAN slicing, and the validity of CUPS was proved by a testbed facilitated by software defined radio and LTE Release 8.

The main contributions of this treatise are summarized as follows. Firstly, upon invoking stochastic geometry, we provide an analytical model for cellular RANs relying on the beneficial feature of CUPS. Secondly, the rate requirements of users in the CP and the UP are defined independently with the aid of queuing theory. Thirdly, the CP and the UP have different channel models, where the UP experiences both LoS and NLoS transmissions, while the CP mainly exhibits NLoS propagation. Finally, since providing reliable coverage is of the highest priority, we analyze the coverage probability for the CP. By contrast, energy efficiency is studied for the UP analysis because of the high-rate data requirement. Our numerical results show that the coverage probability of the CP can be improved upon increasing the service rates or the density of the CBSs, and there exists a TBS density that maximizes the energy efficiency of the UP, while additionally

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improving the area spectral efficiency (ASE). To the best of our knowledge, this is the first contribution jointly modelling and analyzing cellular radio access networks (RANs) relying on CUPS, while considering different rate requirements in LoS/NLoS scenarios in the CP as well as in the UP and invoking stochastic geometry and queuing theory.

II. SYSTEM MODEL AND CHANNEL MODELS

We consider downlink (DL) transmissions in a cellular RAN relying on CUPS, where CBSs and TBSs are deployed following two independent homogeneous spatial Poisson point processes (PPPs), i.e., Ψ_C and Ψ_T , on the plane \mathbb{R}^2 with density of λ_C and λ_T , respectively. The cellular users are distributed following another independent PPP Ψ_U with density λ_u . All users receive control signals from their serving CBSs. The users with low and high user data-rate requirements are connected to their serving CBSs and TBSs, respectively. We denote the probability of a user requiring a high DL data-rate by p_u . The time intervals of each user accessing its serving CBS and TBS obey exponential distributions with their mean value of $1/\lambda_a$ (sec) and $1/\lambda_b$ (sec), respectively. The arrival rates of a user's request at a CBS or a TBS are λ_a and λ_b , respectively. When a user is allocated a channel from its serving CBS and TBS, its channel holding time follows two independent exponential distributions, which have a mean of $1/\mu_c$ (sec) and $1/\mu_t$ (sec), respectively.

A. Traffic Model for the Control Plane

According to the Erlang model[10], the total user traffic load over an area covered by a CBS is $\nu^C = \lambda_a N_C$, where $N_C = \lambda_u/\lambda_C$ is the average number of users served by each CBS. Based on the queuing theory, the queuing model of each CBS is $M/M/C_C/C_C$ associated with an Erlang load of $a_c = \nu^C/\mu_c$, where we assumed that each CBS has C_C channels.

The probability that j channels are being used at a CBS is

$$\phi_j^C = \frac{(a_c)^j}{j!} \phi_0^C = \frac{(a_c)^j}{j!} \left(\sum_{k=0}^{C_C} \frac{(a_c)^k}{k!} \right)^{-1}, \quad (1)$$

where $j \in \{1, \dots, C_C\}$. We denote the probability of a neighboring CBS interfering with a typical CBS by ρ^C [10], where $\rho^C = \sum_{j=1}^{C_C} \frac{j}{C_C} \phi_j^C$. Therefore, the interfering CBSs are distributed following a PPP, which is denoted by Ψ_I^C , with a density of $\lambda_I^C = \rho^C \lambda_C$. The CBSs DL transmission power on each channel is set to P_t^C .

B. Channel Model for the Control Plane

The DL channels of CBSs exhibit both small-scale Rayleigh fading. We assume NLoS transmissions with a path loss exponent of α in the CP due to the larger coverage area. Specifically, the path gain in the CP is given by

$$\Theta^{CP} = Ar^{-\alpha}, \quad (2)$$

where r denotes the transmission distance in the CP. Furthermore, A stands for the corresponding path gain at a reference distance of 1 meter in the CP. We set $A = 1$ for simplicity.

For a user associated with its nearest CBS, the signal-to-interference ratio-plus-noise (SINR) is expressed as

$$\text{SINR}_C = \frac{P_t^C h r^{-\alpha}}{\sum_{k \in \Psi_I^C} P_t^C h_k r_k^{-\alpha} + \sigma^2}, \quad (3)$$

where h and h_k denote the channels' power gain in the desired link and in the interfering link spanning from the k th interfering CBS, which follow two independent exponential distributions with unity mean, respectively. While r and r_k denote the length of the desired link and of the interfering link spanning from the k th interfering CBS, respectively, and σ^2 denotes the additive noise. For CBSs, providing reliable CP coverage has a higher priority than offering high data transmission capacity. Hence, we analyze the coverage probability of the CP, which is defined as the probability that the DL SINR of a user is above a predefined threshold x_0 , i.e.,

$$p^{\text{cov}}(x_0, \lambda_C, \alpha) = \Pr[\text{SINR}_C > x_0]. \quad (4)$$

C. Traffic Model for the User Plane

Considering the potentially high density of TBSs, we use a path loss model to include both LoS and NLoS transmissions depending on the ranges of links. The path gain in the UP is given by

$$\Theta^{UP} = \begin{cases} A^L r^{-\alpha^L} & \text{if LoS w.p. } P_T^L(r) \\ A^N r^{-\alpha^N} & \text{if NLoS w.p. } 1 - P_T^L(r) \end{cases}, \quad (5)$$

where α^L and α^N represent the path loss exponents for LoS and NLoS transmissions, respectively. The corresponding path gains at a reference distance of 1 meter in UP are given by A_L and A_N , respectively. $P_T^L(r)$ is the LoS probability function of TBSs, which is defined as

$$P_T^L(r) = \begin{cases} 1 - \frac{r}{d_0}, & 0 < r \leq d_0 \\ 0, & r > d_0 \end{cases}, \quad (6)$$

where d_0 denotes the reference distance, which is given by $d_0 = 300$ m. The probability of a link with a range of r being LoS or NLoS is $P_T^L(r)$ or $P_T^N(r)$, respectively, which satisfy $P_T^L(r) + P_T^N(r) = 1$.

Note that because CBSs and TBSs transmit their signals in the low and high carrier frequencies, respectively, the path gain of the CP can be much larger than that of the UP. Thus, we set $A = 1$ for simplicity and use the path loss exponents α , α^L and α^N having different values.

If there exists a LoS path between a TBS and a typical user, then the TBS is considered as a LoS TBS. For a typical user, all the LoS TBSs form the set Ψ_L , while the other TBSs are considered to be NLoS TBSs and form the set Ψ_N . We have $\Psi_L \cup \Psi_N = \Psi_T$. Note that Ψ_L and Ψ_N are two independent non-homogeneous PPPs with densities of $\lambda_L(r) = P_T^L(r) \lambda_T$ and $\lambda_N(r) = P_T^N(r) \lambda_T$, respectively. For a typical user in the UP, if its Φ_L is not empty, then the user is served by the LoS TBS that offers the highest received power; otherwise, the user is served by the closest NLoS TBS. The total user traffic over an area covered by a LoS TBS and a NLoS TBS is $\nu^L(r) = \lambda_b N_L(r)$ and $\nu^N(r) = \lambda_b N_N(r)$,

where the average numbers of users served by each LoS TBS and each NLoS TBS are $N_L(r) = p_u \lambda_u / \lambda_L(r)$ and $N_N(r) = p_u \lambda_u / \lambda_N(r)$, respectively. Each TBS has access to a total of C_T channels, and thus a LoS TBS and a NLoS TBS can be modeled as an $M/M/C_T/C_T$ queue with an Erlang load of $a_l(r) = \nu^L(r) / \mu_t$ and $a_n(r) = \nu^N(r) / \mu_t$. The probabilities that $j \in 1, \dots, C_T$ channels are being used in a LoS TBS and a NLoS TBS are $\phi_j^L(r) = \frac{(a_l(r))^j}{j!} \phi_0^L(r)$ and $\phi_j^N(r) = \frac{(a_n(r))^j}{j!} \phi_0^N(r)$, respectively.

Let's denote the probabilities of a neighbor LoS TBS and a neighbor NLoS TBS imposing interference on the typical user by $\rho^L(r)$ and $\rho^N(r)$ [10], respectively. We have $\rho^L(r) = \sum_{j=1}^{C_T} \frac{j}{C_T} \phi_j^L(r)$ and $\rho^N(r) = \sum_{j=1}^{C_T} \frac{j}{C_T} \phi_j^N(r)$. For the typical user, the interfering LoS TBSs and the interfering NLoS TBSs are distributed obeying two independent PPPs denoted by Φ_I^L and Φ_I^N , which have densities of $\lambda_I^L(r) = \rho^L(r) \lambda_L(r)$ and $\lambda_I^N(r) = \rho^N(r) \lambda_N(r)$, respectively. Every channel of a TBS has a DL transmission power P_t^T per channel.

When the number of active users exceeds the number of channels available for a typical TBS, a call will be blocked. The user blocking probabilities of a LoS TBS and of a NLoS TBS are $P_b^L(r) = \phi_{C_T}^L(r)$ and $P_b^N(r) = \phi_{C_T}^N(r)$, respectively.

D. Channel Model for the User Plane

In the UP, a typical user associates with a TBS controlled by the maximum received signal power based user-association scheme. The SINR of a user associated with its serving LoS TBS or its serving NLoS TBS is

$$\text{SINR}_U = \frac{P_t^T A^U h^U r^{-\alpha^U}}{I_T + \sigma^2}, \quad (7)$$

where $U \in \{L, N\}$, A^U is the path loss constant at a reference distance, and α^U is the path loss exponent. The channel's power gain h^U follows an exponential distribution with unity mean; I_T is the cumulative interference power received from all the interfering LoS TBSs and NLoS TBSs, which can be expressed as

$$I_T = \sum_{j \in \Psi_I^L} P_t^T A^L h_j^L r_j^{-\alpha^L} + \sum_{k \in \Psi_I^N} P_t^T A^N h_k^N r_k^{-\alpha^N}. \quad (8)$$

Because the TBSs are responsible for the high speed data transmissions, we analyze the energy efficiency [9-10] of the UP, which is defined by $\eta_E = \frac{\text{ASE}}{\text{PCA}}$.

The ASE is quantified in terms of b/s/Hz/m², while PCA is the transmission power consumption per unit area (PCA) in W/m², both of which will be explained in the next section.

III. PROBLEM FORMULATION AND SOLUTION

In this section, we will derive the coverage of the CP and the energy efficiency of the UP by exploiting the properties of PPP, respectively.

A. Coverage Probability of the Control Plane

Upon substituting (3) into (4), we obtain the coverage probability of the CP as

$$\begin{aligned} p_C^{\text{cov}}(x_0) &= \int_{r>0} \mathbb{P}[\text{SINR}_C > x_0 | r] f_R(r) dr \\ &= \int_{r>0} \exp\left(-\pi r^2 \lambda_I^C Z(x_0, \alpha) - \frac{x_0 \sigma^2 r^\alpha}{P_t^C}\right) f_R(r) dr, \end{aligned} \quad (9)$$

where the probability density function (PDF) of the transmission distance between the user and its serving CBS $f_R(r)$ in (9) is given by $f_R(r) = 2\pi \lambda_C r e^{-\pi \lambda_C r^2}$, $Z(x_0, \alpha) = x_0^{2/\alpha} \int_{x_0^{-2/\alpha}}^{\infty} \frac{1}{1+u^{\alpha/2}} du$. (9) is exploiting the properties of the exponential distribution $h_k \sim \exp(1)$ [13].

B. Energy Efficiency of the User Plane

We define the event that the typical user is connected to a LoS TBS or a NLoS TBS as X^L and X^N with probabilities of H^L and H^N , respectively, i.e., $H^L = \mathbb{P}[X^L]$ and $H^N = \mathbb{P}[X^N]$. We have $H^L + H^N = 1$. The distances between the typical user and its nearest LoS TBS as well as its nearest NLoS TBS are denoted by D_T^L and D_T^N , respectively. Then H^L can be expressed as

$$\begin{aligned} H^L &= \mathbb{P}\left[P_t^T A^L (D_T^L)^{-\alpha^L} > P_t^T A^N (D_T^N)^{-\alpha^N} \mid D_T^L = r\right] \\ &= 2\pi \lambda_T \int_0^\infty \theta^L r P_T^L(r) dr. \end{aligned} \quad (10)$$

Similarly, H^N can be formulated as

$$H^N = 2\pi \lambda_T \int_0^\infty \theta^N r (1 - P_T^L(r)) dr, \quad (11)$$

where $\theta^L = \exp(-\int_0^r \zeta^L(u) du - \int_0^{k_1 r^{\frac{\alpha^L}{\alpha^N}}} \zeta^N(u) du)$, and $\theta^N = \exp(-\int_0^{k_2 r^{\frac{\alpha^N}{\alpha^L}}} \zeta^L(u) du - \int_0^r \zeta^N(u) du)$. $k_1 = (A^L/A^N)^{1/\alpha^N}$, $k_2 = (A^N/A^L)^{1/\alpha^L}$, $\zeta^L(u) = 2\pi u \lambda_T P_T^L(u)$ and $\zeta^N(u) = 2\pi u \lambda_T (1 - P_T^L(u))$.

In the UP, the user will connect to its nearest LoS TBS, namely to the one that provides a high received power instead of the nearest TBS associated with a NLoS path [14]. When the typical user is associated with a LoS TBS or a NLoS TBS, the distance between the user and its serving TBS is denoted as R_T^L and R_T^N , respectively. Furthermore, $R_T^L > r$ and $R_T^N > r$ are equivalent to $D_T^L > r$ and $D_T^N > r$, given that the typical user is connected to a LoS TBS and a NLoS TBS, respectively. Hence the PDFs of the two distances can be expressed as

$$f_{R_T^U}(r) = \frac{d(1 - [R_T^U > r])}{dr} = \frac{2\pi \lambda_T r P_T^U(r) \theta^U}{H^U} \quad (12)$$

where $U \in \{L, N\}$.

According to (9), given a predefined threshold x_0 , the coverage probability of a user served by a LoS TBS is given by

$$p_L^{\text{cov}}(x_0) = \int_{r>0} \exp\left(-\frac{x_0 \sigma^2 r^{\alpha^L}}{P_t^L A^L}\right) \mathcal{L}_{I_T}\left(\frac{x_0 r^{\alpha^L}}{P_t^L A^L}\right) f_{R_T^L}(r) dr, \quad (13)$$

where $f_{R_T^L}(r)$ is provided in (12). According to the properties of PPPs, the Laplace transform of I_T in (8) is obtained as

$$\mathcal{L}_{I_T}(s) = E_{h_k} \left[\exp \left(-s \sum_{k \in \Psi_I^L} P_t^T A^L h_k^L r_k^{-\alpha^L} \right) \right] \times E_{h_j} \left[\exp \left(-s \sum_{j \in \Psi_I^N} P_t^T A^N h_j^N r_j^{-\alpha^N} \right) \right]. \quad (14)$$

Substituting $s = \frac{x_0 r^{\alpha^L}}{P_t^T A^L}$ into (14), we arrive at

$$\mathcal{L}_{I_T} \left(\frac{x_0 r^{\alpha^L}}{P_t^T A^L} \right) = \exp \left(-\pi r^2 \lambda_I^L(r) Z(x_0, \alpha^L) \right) \times \exp \left(-\pi r^2 \lambda_I^N(r) Z \left(\frac{x_0 A^N}{A^L}, \alpha^N \right) \right). \quad (15)$$

Similarly, we obtain the expression of the coverage probability $p_N^{cov}(x_0)$, which is the same as $p_L^{cov}(x_0)$, except that λ_L , A^L and α^L are replaced by λ_N , A^N and α^N .

Upon substituting $s = \frac{x_0 r^{\alpha^N}}{P_t^T A^N}$ into (14), we have

$$\mathcal{L}_{I_T} \left(\frac{x_0 r^{\alpha^N}}{P_t^T A^N} \right) = \exp \left(-\pi r^2 \lambda_I^L(r) Z \left(\frac{x_0 A^L}{A^N}, \alpha^L \right) \right) \times \exp \left(-\pi r^2 \lambda_I^N(r) Z(x_0, \alpha^N) \right). \quad (16)$$

In (15) and (16), $Z(x_0, \alpha^L)$, $Z \left(\frac{x_0 A^N}{A^L}, \alpha^N \right)$, $Z \left(\frac{x_0 A^L}{A^N}, \alpha^L \right)$ and $Z(x_0, \alpha^N)$ follow the expression of $Z(x_0, \alpha)$, as defined for (9).

We denote the area spectral efficiencies provided by LoS TBSs and NLoS TBSs, respectively, as η^L and η^N , which are expressed as

$$\eta^U = \mathbb{E} [\log_2 (1 + SINR_U)] = \frac{1}{\ln 2} \int_0^\infty \frac{p_U^{cov}(x_0)}{x_0 + 1} dx_0. \quad (17)$$

We denote the expectations of user blocking probabilities by P_b^L and P_b^N , which represent the average probability of a typical user's call being blocked by a LoS TBS or a NLoS TBS, respectively. Then P_b^L and P_b^N can be calculated as

$$P_b^L = \int_0^\infty P_b^L(r) f_{R_T^L}(r) dr, \quad (18)$$

$$P_b^N = \int_0^\infty P_b^N(r) f_{R_T^N}(r) dr. \quad (19)$$

For a given TBS intensity λ_T , the ASE in the UP is expressed as follows

$$ASE = p_u \lambda_u (1 - P_b^L) H^L \eta^L + p_u \lambda_u (1 - P_b^N) H^N \eta^N. \quad (20)$$

The PCA in the UP can be expressed as

$$PCA = \int_0^\infty (\lambda_L(r) \varpi^L(r) f_{R_T^L}(r) + \lambda_N(r) \varpi^N(r) f_{R_T^N}(r)) dr, \quad (21)$$

where $\varpi^L(r)$ and $\varpi^N(r)$ denote the power consumptions of a LoS TBS and a NLoS TBS, respectively, which are given by

$$\varpi^U(r) = \sum_{j=1}^{C^T} j P_t^T \phi_j^U(r) + P_{OM} \phi_0^U(r), \quad (22)$$

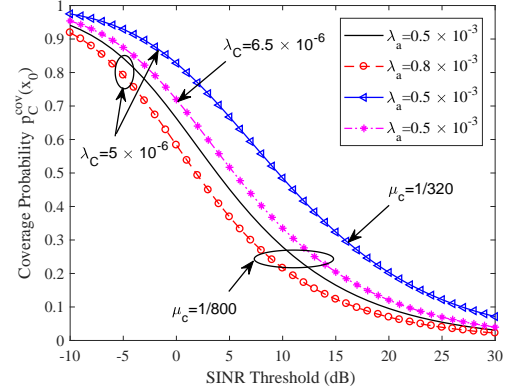


Fig. 1. The coverage probability of the CP vs. the SINR threshold λ_T .

where $U \in \{L, N\}$ and P_{OM} is the power assigned to cooling and other mains-operated functions. If no user requests any service within a TBS's coverage, no channel will be occupied. Therefore, the TBS remains sleeping state consuming a power of P_{OM} .

Combining (20) and (21), we obtain the energy efficiency of the UP as

$$\eta_E = \frac{p_u \lambda_u (1 - P_b^L) H^L \eta^L + p_u \lambda_u (1 - P_b^N) H^N \eta^N}{\int_0^\infty [\lambda_L(r) \varpi^L(r) f_{R_T^L}(r) + \lambda_N(r) \varpi^N(r) f_{R_T^N}(r) dr] dr}. \quad (23)$$

IV. NUMERICAL RESULTS AND DISCUSSIONS

In this section, we present our numerical results to show the performance of cellular RANs relying on CUPS in accordance with our theoretical analysis. According to [16] - [18], we use the parameters $P_t^C = 1$ W, $P_{t1}^T = 0.4$ W, $P_{OM} = 0.25$ W, $\sigma^2 = -90$ dBm, $A^L = 10^{-4.11}$, $A^N = 10^{-3.01}$, $\alpha = 4$, $\alpha^L = 2.5$, $\alpha^N = 3.5$, $C_C = 50$, $C_T = 5$, $p_u = 0.3$, $d_0 = 300$, $\lambda_a \in \{0.5 \times 10^{-3}, 0.8 \times 10^{-3}\}$, $\lambda_b \in \{4 \times 10^{-3}, 5 \times 10^{-3}\}$, $1/\mu_c \in \{800, 320\}$ (sec), and $1/\mu_t \in \{80, 64\}$ (sec), unless specified otherwise.

Fig. 1 shows the coverage probability of the CP versus the SINR threshold x_0 , for various λ_a , μ_c and λ_C values. Note that the coverage probability increases with increasing the service rate μ_c and λ_C , or with reducing either the access rate λ_a or the SINR threshold x_0 . The reduction of λ_a , λ_C and μ_c leads to an increase in λ_I^C , causing a decrease in the coverage probability. Therefore, when the arrival rate of user requests is high, a higher coverage probability can be achieved by increasing the CBSs' deployment density or by increasing each CBS's service rate.

As illustrated in Fig. 2, when the TBSs are sparsely deployed (from 0.5×10^{-5} to 1.2×10^{-5} BSs/m²) and thus the interference is limited, the ASE of the UP increases quickly due to the reduction of the transmission distance and path loss. Upon the further increasing the density from 1.2×10^{-5} to 2×10^{-5} BSs/m², the ASE shows a slight downward trends. Because the average radius of a TBS coverage area ranges from about 126 m to about 162 m in this density region, a user is connected to a LoS TBS with a high probability. The path loss model of interference is changed from NLoS

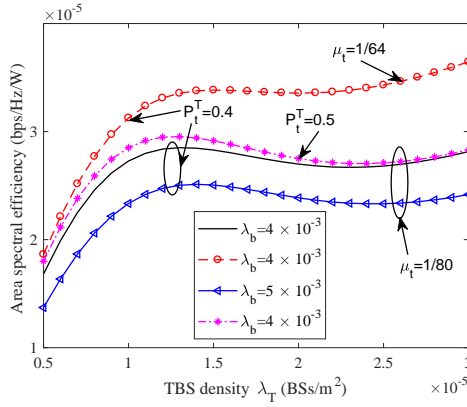


Fig. 2. The ASE of the user plane vs. the TBSs' density λ_T .

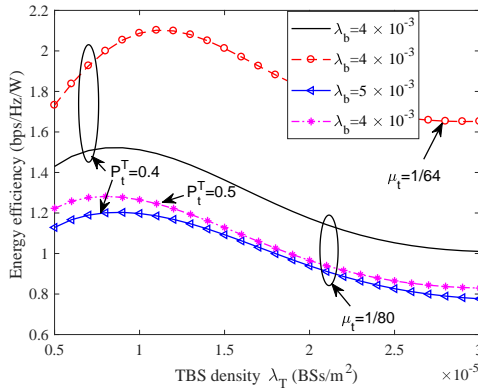


Fig. 3. The energy efficiency of the UP vs. the TBSs' density λ_T .

to LoS, thereby increasing the interference imposed on the user. When the TBSs are deployed more densely, the coverage probability tends to a constant upon increasing λ_T (from 2×10^{-5} to 3×10^{-5} BSs/m²), because increasing the LoS TBS power is counter-balanced by the increased interference power. Therefore, the ASE shows an increasing trend. Similar trends can be found in [19] and [20].

Fig. 3 illustrates the energy efficiency of the UP versus the TBS density λ_T for various λ_b , μ_t and P_t^T values. Intuitively, there exists a TBS density λ_T^* that maximizes the energy efficiency. More specifically, if $\lambda_T < \lambda_T^*$, the energy efficiency increases with respect to λ_T . By contrast, if $\lambda_T > \lambda_T^*$, the energy efficiency decreases slowly with λ_T . Observe from Fig.1(b) and Fig.1(c), that the UP cannot simultaneously attain the highest ASE and the highest energy efficiency at the same density λ_T . The maximum energy efficiency of the UP can be achieved by optimizing the TBS density. When the arrival rate λ_b is high, it is necessary to appropriately reduce P_t^T or increase μ_t in order to increase the energy efficiency.

V. CONCLUSIONS

In this paper, we have analyzed the performance of cellular RANs relying on CUPS, where the CP and the UP have different dynamic traffic load requirements and channel models. By relying on a sophisticated combination of stochastic geometry and queuing theory, we derive the coverage probability of the

CP and the energy efficiency of the UP. Interestingly, base on numerical results, we have found that increasing the CBSs' deployment density or service rates increases the coverage probability of the CP, but there is a performance-cost trade-off. Although the ASE and the energy efficiency of the UP cannot be simultaneously maximized by optimizing the density of TBS, the maximum energy efficiency of the UP can be achieved by optimizing the density of TBSs, while achieving a high ASE.

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